Towards Multimedia Fragmentation

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Abstract. Database fragmentation is a process for reducing irrelevant data accesses by grouping data frequently accessed together in dedicated segments. In this paper, we address multimedia database fragmentation by extending existing fragmentation algorithms to take into account key characteristics of multimedia objects. We particularly discuss multimedia primary horizontal fragmentation and provide a partitioning strategy based on low-level multimedia features. Our approach particularly emphasizes the importance of multimedia predicates implications in optimizing multimedia fragments. To validate our approach, we have implemented a prototype computing multimedia predicates implications. Experimental results are satisfactory.

Keywords: Multimedia fragmentation, Range and KNN operators, predicates implication, objects classification

1 Introduction

Since the last two decades, multimedia data are of key importance in many application areas such as medicine, surveillance, cartography, meteorology, security, visual data communications, etc. Hence, the need for systems that can catalog, store, and efficiently retrieve relevant distributed multimedia data is becoming very high. Initially, research in multimedia management has been handled separately by database management and computer vision communities. As a result, different types of features have been used, in the literature, for multimedia data management. Low-level features such as color, texture, shape, layout, etc. are used by the computer vision research community, while meta-data and semantic based features are widely used by the database management community to describe data context and semantics. Emerging applications in distributed environments create an increasing demand on the performance of multimedia systems, requiring new data partitioning techniques to achieve high resource utilization and increased concurrency and parallelism. Several continuing studies are aimed at building distributed multimedia databases management systems MMDBMS [20]. Nevertheless, most existing systems lack a formal framework to adequately provide full-fledge multimedia operations.

Traditionally, partitioning techniques are used in distributed system design to reduce accesses to irrelevant data. Three main fragmentation techniques have been defined for relational databases: horizontal fragmentation HF, vertical fragmentation (VF), and hybrid or mixed fragmentation (MF). These techniques have been recently extended for object oriented databases. However, multimedia data fragmentation issues haven't been addressed in current systems.

Multimedia fragmentation is a relatively complicated issue owing to the complexity of the multimedia data itself; different multimedia data types (video, audio, image and/or text), frequently used with various formats, as well as the intricacy of the description of physical and/or semantic multimedia data. In this paper, we address primary horizontal fragmentation in distributed multimedia databases and analyze the impact of multimedia operators and predicates. We particularly address multimedia predicates implication required in current fragmentation algorithms such as Make_Partition and Com_Min [2, 11, 12]. We also present our prototype with corresponding experimental results conducted to validate our approach.

The remainder of this paper is organized as follows. Section 2 briefly reviews background in DB fragmentation. Section 3 presents a motivation example. Section 4 details our multimedia fragmentation process. Section 5 presents our prototype and experimental tests. Finally, section 6 concludes and draws future directions.

2 Background

Fragmentation techniques for distributed DB systems aim to achieve high resource utilization and performance [5]. This is addressed by removing irrelevant data accessed by applications and by reducing data exchange among sites [1]. In this section, we briefly present traditional database fragmentation approaches, depicting the evolution from relational to object oriented DBMS, and focus on horizontal fragmentation algorithms. In essence, there are three fundamental fragmentation strategies: Horizontal Fragmentation (HF), Vertical Fragmentation (VF) and Mixed Fragmentation (MF).

HF underlines the partitioning of an entity/class in segments of tuples/objects verifying certain criteria. The generated horizontal fragments have the same structure as the original entity/class. Horizontal fragmentation is generally categorized in two types: Primary HF and Derived HF. PHF is the partitioning of an entity based on its attributes' values [12]. DHF denotes the partitioning of an entity (called member) based on links with other entities (called owners) [12]. In other words, it is the partitioning of an entity/class in terms of the PHF of another entity/class [1] taking into consideration their inner-links.

VF breaks down the logical structure of an entity/class by distributing its attributes/methods over vertical fragments, which would contain the same tuples/objects with different attributes [1]. The unique tuple/object identifier (id) is kept in all vertical fragments [7] so that the DBMS can link related segments.

MF is a hybrid partitioning technique where horizontal and vertical fragmentations are simultaneously applied on an entity/class [11].

To the best of our knowledge, two main algorithms for the PHF of relational DBMS are provided in the literature: *Com_Min* developed by Oszu and Valduriez [12] and

Make_Partition Graphical Algorithm developed by Navathe *et al.* [10] (used essentially for vertical fragmentation). The *Com_Min* algorithm generates, from a set of simple predicates applied to a certain entity, a complete and minimal set of predicates used to determine the minterm fragments corresponding to that entity. A minterm is a conjunction of simple predicates [2] associated to a fragment. *Make_Partition* generates minterm fragments by grouping predicates having high affinity towards one another. The number of minterm fragments generated by *Make_Partition* is relatively smaller than the number of *Com_Min* minterm fragments [15] (the number of minterm fragments generated by *Com-Min* being exponential to the number of simple predicates considered).

Similarly, there are two main algorithms for the PHF of object oriented DBMS: one developed by Ezeife and Barker [6] using Com_Min [12], and the other developed by Bellatreche *et al.* [2] on the basis of Make_Partition [10]. The use of Com_Min or Make_Partition is the major difference between them.

3 Motivation

In order to use current partitioning approaches, widely employed in traditional databases, for fragmenting multimedia data, several issues should be studied and extended. On one hand, to achieve fragmentation, current algorithms require as an input parameter [6] the database conceptual schema (CS). This requirement is not always fulfilled in some multimedia databases due to the unstructured (or semi-structured) and complex nature of multimedia data. On the other hand, multimedia queries contain new operators handling low-level and semantic features. These new operators should be considered when studying predicates and particularly predicate implications. For example, let us consider the following predicates used to search for photos similar to given photos in an Employee multimedia database as shown below.

Predicate	P1	P2	P3	P4 Emp_photo			
Attribute	Emp_photo	Emp_photo	Emp_photo				
Operator ¹	$Range_Sim_{\varepsilon l}$	$Range_Sim_{\epsilon^2}$	$Range_Sim_{\epsilon^3}$	KNN			
Value			PROPOSIL	I PROPOSILI			
Parameter	$\varepsilon_2 > \varepsilon_1$	$\epsilon_2 > \epsilon_1$	$\varepsilon_3 > \varepsilon_1$	K=3			

In current approaches, the following predicates are considered different and analyzed separately:

- P₁ and P₂: two range queries with different parameters (radius)
- P₁ and P₃: two range queries with different parameters and values
- P₃ and P₄: two different operators

However, in multimedia applications, P_1 would also retrieve objects belonging to results of queries based on P_2 and P_3 . Likewise, P_4 may return a subset of P_3 's results. Thus, we can say that P_2 and P_3 infer P_1 (denoted by $P_1 \rightarrow P_2, P_3$), and consider only the results returned by P_2/P_3 , thus eliminating P_1 .

¹ More details about multimedia operators will be given later.

It is important to notice that ignoring such implications between predicates can lead, in multimedia applications, to higher computation costs when creating fragments, bigger fragments which is very restrictive for multimedia storage, migration, and retrieval, as well as data duplication on several sites. In [2, 11], the authors have only highlighted the implication issue importance, but have not well detailed nor identified the various kinds of implications. These issues will be tackled in following paragraphs.

4 Multimedia Primary Horizontal Fragmentation

In this section, we start by introducing some concepts and definitions necessary to tackle multimedia primary horizontal fragmentation. We develop subsequently additional steps to be integrated in current approaches, allowing adequate multimedia data fragmentation processing.

4.1 Definitions

4.1.1 Multimedia Object

A multimedia object is described by a set of attributes, related to a set of meta-data. It can be formally depicted as a set of attribute (a_i) and value (v_i) doublets: O { (a_1, v_1) ; (a_2, v_2) , ..., (a_n, v_n) }. Multimedia attributes and values can be *simple* (like color = "red"), *complex* (color histogram, texture, shape, etc.) or the raw data (BLOB files) of multimedia objects.

4.1.2 Multimedia Type

A multimedia type allocates a set of attributes used to describe multimedia objects corresponding to that type: $T(a_1, a_2, a_3, \dots, a_n)$. We consider that two objects, described by the same attributes, are of the same type.

4.1.3 Multimedia Query

A multimedia query is written as follows [2, 9]:

 $q = \{(Target clause), (Range clause), (Qualification clause)\}, where:$

- *Target clause*: contains multimedia attributes returned by the query
 - **Range clause:** gathers the entities (tables/lasses) accessed by the query, to which belong *target clause* and *qualification clause* attributes
 - **Qualification clause:** is the query restriction condition, a Boolean combination of predicates, linked by logical connectives Λ , v, \neg

4.1.4 Multimedia Operators and Predicates

As mentioned before, multimedia information introduces new types of data and new operators and predicates. In the following, we explain multimedia operators and predicates related to low-level features. Note that semantic similarity operators are out of this paper's scope and will be detailed in future studies.

4.1.4.1 Multimedia Operators

In multimedia databases, objects are widely described using vector spaces with numeric attributes, such as shape or color descriptors. Thus, in order to retrieve multimedia data, dedicated similarity queries are used, involving *range queries* and/or *k-nearest neighborhood* operators. Formal definitions are given thereafter.

4.1.4.1.1 Multimedia Range Query Operator

A range query operator $\overline{\theta}$ returns the set of objects V_j of an object value V_i located within a certain range ε from V_i using a distance function *D* (cfr. *Figure 1*). It can be formally written as:

Range Query(V_i,
$$\overline{\theta}$$
, ε) = N ^{ε} _{$\overline{\alpha}$} (V_i) = {V_j / D(V_i, V_j) $\leq \varepsilon / \varepsilon \in \mathbb{R}$ (1)

The function D can be the classic Euclidean distance, a weighted Euclidean distance, a quadratic form distance, etc.



Fig. 1. Visualizations of a range query operator θ

A range query operator $\overline{\theta}$ has the following interesting properties, useful for optimizing the computation process:

- $N_{\overline{\theta}}^{\varepsilon_{i}}(V_{i}) \subseteq N_{\overline{\theta}}^{\varepsilon_{j}}(V_{i}) \text{ if } \varepsilon_{i} \leq \varepsilon_{j}$
- if $N_{\overline{\theta}}^{\varepsilon_i}(V_i) \subseteq N_{\overline{\theta}}^{\varepsilon_j}(V_j)$ and $N_{\overline{\theta}}^{\varepsilon_j}(V_j) \subseteq N_{\overline{\theta}}^{\varepsilon_i}(V_l) \to N_{\overline{\theta}}^{\varepsilon_i}(V_i) \subseteq N_{\overline{\theta}}^{\varepsilon_i}(V_l) \ \forall \varepsilon_i, \varepsilon_j, \varepsilon_l$

4.1.4.1.2 Multimedia KNN Operator

A K-Nearest Neighborhood (KNN) operator θ returns the set of K neighbors of an object value V_i located into either a ranged or unlimited domain space, using a distance *D* [3, 20]. It could be formally written as follows:

$$KNN(V_{i}, \overline{\theta}, k)_{\varepsilon} = N_{\overline{\theta}}^{k}(V_{i})_{\varepsilon} = \left\{ V_{j=1..k} / D(V_{i}, V_{j}) \leq D(V_{i}, V) \right\}$$

$$\forall V \notin N_{\overline{\theta}}^{k}(V_{i}), \text{ where } k \in \mathbb{N} \text{ and } Max(D(V_{i}, V)) \leq \varepsilon / \varepsilon \in \mathbb{R}^{*} \cup \{ \bot \}$$

If $\varepsilon = \bot$, the domain space is unlimited (2)

As for range query operators, a KNN operator can be observed as a visual object in function of values dimensions. *Fig.* 2 shows a ranged 2D KNN operator with k=3.



Fig. 2. Visualizations of a ranged 2D KNN operator θ

A KNN operator $\vec{\theta}$ has the following properties:

- $N_{\bar{\theta}}^{k_i}(V_i) \subseteq N_{\bar{\theta}}^{k_j}(V_i) \text{ if } k_i \leq k_j$
- if $N_{\hat{a}}^{k_i}(V_i) \subseteq N_{\hat{a}}^{k_j}(V_j)$ and $N_{\hat{a}}^{k_j}(V_j) \subseteq N_{\hat{a}}^{k_i}(V_l) \to N_{\hat{a}}^{k_i}(V_i) \subseteq N_{\hat{a}}^{k_i}(V_l) \ \forall k_i, k_j, k_l$

4.1.4.2 Multimedia Predicates

A multimedia predicate \hat{P} is defined as follows:

$$\hat{\mathbf{P}}_{i} = (\mathbf{A}_{i} \ \theta_{m} \ \mathbf{V}_{i})$$

Where:

- A_i is a multimedia attribute or object
- V_i is a value in the domain of A_i or a multimedia object
- $\theta_m = \theta_t \cup \{\bar{\theta}, \bar{\theta}\}$ where θ_t is a traditional operator such as a comparison

operator (=, <, \leq , >, \geq , \neq), or a set operator (contained-in, set-equality, ...), etc.

4.2 Steps for Multimedia Data Primary Horizontal Fragmentation

Before applying current fragmentation approaches, several steps should be executed in order to support and provide relevant multimedia data fragmentation. We suggest integrating the steps detailed below.

```
Multimedia_fragmentation_pre-processing ()
```

```
Begin

Multimedia_Types_Classification()

For each multimedia Type

Predicates_Grouping()

Multimedia_Predicates_implication()

EndFor

End
```

4.2.1 Classification of Multimedia Objects

By applying existing horizontal fragmentation algorithms to a multimedia database, we attain non consistent horizontal fragmentation criteria (minterms). Suppose that *Camera Position, Audio Frequency* and *Dominant Color* are three multimedia attributes describing Video, Audio and Image objects respectively. The following Boolean expression: *CameraPosition* = "North West" A Audio Frequency = "6 KHz" A DominantColor = ((10; 10; 10), RGB) is a non consistent minterm, specifying

criteria on "heterogeneous" attributes describing multimedia objects of different types, therefore producing an empty horizontal fragment.

In order to attain coherent minterms, we need to gather related objects together. As mentioned before, we assume that multimedia objects having the same attributes are considered of the same type. The algorithm provided below is used for classifying objects, according to their corresponding types.

Multimedia_Types_Classification ()

Input : MM Output : T_M	<pre>// multimedia objects //set of multimedia types correspond</pre>	ing to objects in MM
Begin		
	For each $Mo_i \in MM$	
	If Mo _i A ≠ all T _i A	// Adding a new type corresponding to the object Moi
	New T_{n+1}/T_{n+1} . A = Mo _i . A	// if the type isn't considered yet in MM
	$T_{n+1} = T_{n+1} U Mo$	
	Else	
	$T_i = T_i U Mo_i / Mo_i A = T_i A$	// Adding the object Moi to its corresponding type
	Endlf	// if the type is already identified
	Endfor	5. S
End		

4.2.2 Predicate Grouping

It is also important to gather predicates into groups on the basis of operators. Using the algorithm below, two predicate groups are identified: multimedia and traditional. This separation will allow defining appropriate methods for multimedia implication:

$$\mathbf{P}_{i} \xrightarrow{\theta_{m}} \mathbf{P}_{j} \Leftrightarrow \begin{bmatrix} \widehat{\mathbf{P}_{i}} \xrightarrow{\theta} \widehat{\mathbf{P}_{j}} \\ \mathbf{P}_{i} \xrightarrow{\theta} \mathbf{P}_{j} \end{bmatrix} \xrightarrow{\theta} \text{ denotes a multimedia similarity implication}$$

$$\xrightarrow{\theta} \text{ denotes a traditional implication}$$

Recall that traditional implication is out of this paper's scope.

Predicates_grouping ()

Input: Q //set of all user queries Ti //a multimedia type Output: P,i //a query predicate defined on type T \hat{P}_i //set of multimedia predicates applied on T Ρ, //set of traditional predicates applied on T For each query $Q_i \in Q$ For each $P_i^{i} \in Q_i$ If ($\underline{P}_i^{i} \in \underline{P})$ then $\hat{P_i} = \hat{P_i} \cup P_j^i$ Else $P_i = P_i \cup P_j^i$ Endif EndFor EndFor End

Begin

4.2.3 Multimedia Predicates Implication

Finding inference or implication between predicates is crucial to cutback the number of predicates involved in the fragmentation process [4, 11] (a large number of unnecessary fragments would notionally achieve low system performance). When a predicate P_i implies a predicate P_j (denoted by $P_i \rightarrow P_j$), P_i can be removed from the minterm fragment to which it belongs and replaced by P_j . Predicate implication is taken into consideration in traditional algorithms, mainly in Com_Min [12] and Make_Partition algorithms [10]. In the following, we detail the rules that can be used to determine implication between low-level feature-based predicates, by using both: range query and KNN methods.

4.2.3.1 Range Query Predicates Implication

Two range query predicates $\overline{P_i}$ and $\overline{P_i}$ are in implication if:



Fig. 3. 2D Range Query Predicates Implication

However, if $\varepsilon_i = \varepsilon_j$ and $D(V_i, V_j) \neq 0$ or if $\varepsilon_i - \varepsilon_j < D(V_i, V_j) \le \varepsilon_i + \varepsilon_j$, then there is an intersection between $\overline{P_i}$ and $\overline{P_j}$. Therefore, $\overline{P_i}$ and $\overline{P_j}$ cannot be associated via implication.

4.2.3.2 KNN Predicates Implication

The KNN implications for ranged or unlimited domain space are identical and can only be computed as follows:



Fig. 4. KNN Predicates implication with identical values

Note that two KNN predicates \vec{P}_i and \vec{P}_j identified within two limited ranges ε_i and ε_i are not in implication (like for range queries) if:

$$\{0 < D(V_i, V_j) \le \varepsilon_i - \varepsilon_j \text{ where } \varepsilon_i \text{ and } \varepsilon_j \in [0, 1]\}$$

4.2.3.3 Multimedia Predicates Implication

Using the same reasoning, we consider that two multimedia predicates \hat{P}_i and \hat{P}_j are in implication if:

$$\widehat{P}_{j} \to \widehat{P}_{i} \Leftrightarrow \begin{cases} 0 \le D(V_{i}, V_{j}) \le \varepsilon_{i} - \varepsilon_{j} \text{ and } \left(\widehat{P}_{i} \lor \widehat{P}_{j} \in \left\{\overline{P}\right\} \text{ and } (\overline{\varepsilon} < \overline{\varepsilon}) \right) \\ \text{OR} \\ V_{i} = V_{j} \text{ and } k_{i} \ge k_{j} \end{cases}$$

The first condition allows computing the implication between either two range query predicates or a range query predicate and a ranged KNN predicate. $\vec{\varepsilon}$ is used to designate the range of KNN predicate, and $\vec{\varepsilon}$ to designate the radius of the range predicate. The second condition highlights KNN predicates implication.

The following algorithm generates sets of multimedia predicate implications, IS_i , corresponding to each multimedia type T_i . Note that every set element consists of a doublet of predicates (P_i , P_j), meaning that P_i implies P_j .

Multimedia_Predicates_Implication ()

```
P.
Input:
                                     //set of M multimedia predicates applied on a multimedia type T
                   IS,
Output:
                                     //set of multimedia predicates implications applied on a type T<sub>i</sub>
Variable: P.
                                     //a query predicate defined on type T
Begin
       For each P_i^{i} \in \hat{P_i}
              If j≤M-1 then
                      For each P_{i+1}^{i} \in \hat{P_{i}}
                             If (A_i = A_{i+1}) then
                                                                                                                                                    //same attribute
                                         If (P_i^i.operator = \overline{\theta} \text{ and}(P_{i+1}^i.operator = \overline{\theta} \text{ or } P_{i+1}^i.operator = \overline{\theta})) then
                                                                                                                                                     //R_{j}R_{j+1}, R_{j}K_{j+1}
//P_{j+1}^{i} \rightarrow P_{j}^{i}
                                                    If (\epsilon_j > \epsilon_{j+1}) then
                                                                If 0 \le D(V_j, V_{j+1}) \le \varepsilon_j - \varepsilon_{j+1} then
                                                                          IS_i = IS_i \cup (P_{i+1}^i, P_i^i)
                                                                Endif
                                                    Elseif (\epsilon_{j+1} > \epsilon_j \text{ and } P_{j+1}^i \text{.operator} = \overline{\theta}) then
                                                                                                                                           // R_j R_{j+1}
                                                                                                                                                    ^{\prime\prime}P^i_j \rightarrow P^i_{j+1}
                                                                If 0 \leq D(V_{j+1},V_j) \leq \epsilon_{j+1} - \epsilon_j then
                                                                           IS_i = IS_i \cup (P_i^i, P_{i+1}^i)
                                                                Endif
                                                     Endif
                                         Elseif (P_i^i.operator = \vec{\theta} and P_{i+1}^i.operator = \vec{\theta}) then
                                                                                                                                                         // K<sub>i</sub> K<sub>i+1</sub>
                                         If D(V_i, V_i) = 0 or V_i = V_j then
                                                    If (k_j \ge k_{j+1}) then
```

```
IS_i = IS_i \bigcup (P_{i+1}^i, P_j^i)
                                                   \text{Elseif} \ (k_{j+1} \ \geq \ k_j) \ \text{then}
                                                              IS_i = IS_i \bigcup (P_i^i, P_{i+1}^i)
                                                   Endif
                                        Endif
                             Elseif(P_i^i.operator = \vec{\theta} and P_{i+1}^i.operator = \vec{\theta}) then
                                                                                                                                                      // K<sub>j</sub> R<sub>j+1</sub>
                                        If (\varepsilon_{j+1} > \varepsilon_j) then
                                                                                                                                                      ^{//}P_{i}^{i} \rightarrow P_{i+1}^{i}
                                                   If 0 \leq D(V_{j+1}, V_j) \leq \epsilon_{j+1} - \epsilon_j then
                                                              IS_i = IS_i \cup (P_i^i, P_{i+1}^i)
                                                   Endif
                                       Endif
                            Endif
                     Endfor
              Endif
       Endfor
       IS_i = Optimize(IS_i)
End
Optimize(IS<sub>i</sub>)
  Input: ISi
                                                                                 // set of multimedia predicates implications applied on a type T
      Begin
             For each (P_i^i, P_k^i) \in IS_i
                        For each (P_k^i, P_l^i) \in IS
                                  If (P_i^i \rightarrow P_k^i \text{ and } P_k^i \rightarrow P_l^i) then
                                                IS_i = IS_i \cup (P_i^i, P_i^i)
                                     Endit
                         EndFor
              EndFor
    End
```

4.2.4 Algorithm Complexity

The complexity calculations are carried out below on the basis of the worst case analysis. Suppose n_f represents the largest number of possible fragments, n_o represents the largest number of multimedia objects in a type or a fragment, n_q the largest number of user queries, n_t the largest number of types, n_p the largest number of multimedia predicates, n_i the largest cardinality of the sets IS_i, n_v the largest feature vector dimension involved. Our fragmentation pre-processing algorithm is of time complexity of $O(n_t \times (n_o + n_q \times n_p + n_v \times n_p^2 + n_i^2))$, which simplifies to $O(n_t \times (n_v \times n_p^2))$. Note that the polynomial (quadratic) nature of our features implication computation algorithm ($O(n_v \times n_p^2)$) dominates the complexity formulae and is experimentally demonstrated in our simulation prototype.

4.2.5 Computation Example

In the following, multimedia predicates (range query and KNN) will be illustrated in the same manner for the sake of simplicity:

 $\overline{P} = A \operatorname{Similar}(\varepsilon) V$ and $\overline{P} = A \operatorname{Similar}(k, \varepsilon) V$ where:

- A is a multimedia attribute. In the present example, A stands for Dominant Color : DC
- Similar represents $\overline{\theta}$, the range similarity operator, when the number between brackets ε denotes a real value such as $0.0 \le \varepsilon \le 1.0$; ε designating the similarity range
- Similar stands for $\vec{\theta}$, the KNN operator, when the number between brackets k denotes an integer value ; k representing the number of neighboring objects to be returned by the KNN predicate within a range ϵ

Figure 5 shows three images a, b and c characterized by their feature vector values V_a , V_b and V_c respectively ; V designating, for each image, its *Dominant Color* feature in RGB color space (vector dimension = 3).



Fig. 5. Sample images

We also consider the following two range query predicates:

- **P**₁: DC Similar(ε_1) V₁ and **P**₂: DC Similar(ε_2) V₂ (*DC: Dominant Color*) where V₁ = (22; 22; 22), V₂ = (90; 10; 10), ε_1 = 0.6, and ε_2 = 0.2

Please note that in our similarity computations, we used the following weighted Euclidean distance function:

$$Dist(X,Y) = \frac{\sqrt{\sum_{i=1}^{N} (x_i - y_i)^2}}{\sum_{i=1}^{N} (x_i + y_i)} \in [0,1]$$

N = Max (dim(X), dim(Y)), dim(X) and dim(Y) being the dimensions of vectors X and Y respectively.

Following our multimedia implication computation rules, predicate p_2 implies predicate p_1 ($0 \le \text{Dist}(V_1, V_2) \le \varepsilon_1 - \varepsilon_2$) where:

- $Dist(V_1, V_2) = ((22-90)^2 + (22-10)^2 + (22-10)^2)^{1/2} / (22+90+22+10+22+10) = 0.397$ - and $\varepsilon_1 - \varepsilon_2 = 0.6 - 0.2 = 0.4$

A query utilizing predicate P₁ would return still regions a and b

- Dist(V₁, V_a) = 0.024 ($< \varepsilon_1$, returned object)
- $Dist(V_1, V_b) = 0.399 (< \epsilon_1, returned object)$
- $Dist(V_1, V_c) = 0.662 (> \epsilon_1)$

Whereas a query invoking predicate P2 would return still region b

- $Dist(V_2, V_a) = 0.417 (> \varepsilon_2)$
- Dist(V₂, V_b) = 0.102 (< ε_2 , returned object)
- $\text{Dist}(V_2, V_c) = 0.401 \ (> \varepsilon_2)$

One can clearly realize that the set of multimedia objects returned by P_1 ({a, b}) includes those returned by of P_2 ({b}). If taken into account, such implications would reduce fragment creation computation cost, fragment size and multimedia data duplication on multiple sites.

5 Prototype

To validate our approach, we have implemented a C# prototype called "Multimedia Implication Identifier" encompassing:

- A relational database, storing multimedia objects via Oracle 9i DBMS, described following the multimedia meta-model M² (MPEG-7 compatible) developed by Chalhoub *et al.* in [4].
- A set of interfaces allowing users to formulate simple and complex multimedia queries, providing the ability to select multimedia information.
- Containers for storing user queries, enabling, via specific processes, the computation of query access frequencies which are basically used in the predicate affinity calculations.
- Specific containers undertaking the storage of predicates, utilized by dedicated procedures to calculate predicate implications.

The prototype accepts, as input, multimedia queries. Automatic processes subsequently calculate query access frequencies, identify corresponding predicates, and compute for each multimedia type (represented by a table) its Predicate Usage Matrix (PUM)¹ and its Predicate Affinity Matrix (PAM)² used to measure the affinity between predicates, the PAM taking into account our predicate implication steps.

Note that we chose to present multimedia implications in PAM matrixes, proposed by [15, 4], for the sake of clearness (PAMs being suitable structures for displaying predicate implications). Nevertheless, our algorithm is generic in the sense that it could be equally used with other primary horizontal fragmentation approaches, Com_Min [16] in particular.

5. 1. Simulation example

Among the various tests that were conducted, we present a simple simulation example comparing predicate affinities (PAM) obtained with and without the inclusion of our multimedia physical implication rules. In the following example, multimedia type "Still Region", designating motionless images, is selected for PUM and PAM calculations. Let $Q = \{q_{i = 0 \text{ to } 5}\}$ be a set of user queries defined on "Still Region" Type. Recall that we represent queries following paragraph 4.1.3.

- **q₀:** { (MO); (StillRegion); (ObNature = "vehicule" Λ
- DC Similar(0.3) ((12; 10; 13), (14; 15; 16), (20; 20; 20))) } q1: { (MO); (StillRegion); (ObNature = "vehicule" \land ObColor = "red" \land
- DC Similar(0.2) ((12; 10; 13), (14; 15; 16), (20; 20; 20))) }
- **q₃:** { (MO); (StillRegion); (ObNature = "vehicule" Λ
- DC Similar(3) ((12; 10; 13), (14; 15; 16), (20; 20; 20))) } q_4 : { (MO); (StillRegion); (ObNature = "vehicule" \land ObColor = "red" \land
- DC Similar(1) ((12; 10; 13), (14; 15; 16), (20; 20; 20))) } **qs:** { (MO); (StillRegion); (ObNature = "truck" A ObColor = "red" A
- DC Similar(1) ((9; 8; 7), (7; 8; 7), (10; 11; 10))) }

¹ It contains the predicates used by each query as well as query access frequencies and is subsequently used as input to the PHF process adopted by [11, 2]

² Following [15, 4], the PAM is a square and symmetric matrix where each value $aff(P_i, P_j)$ can be numerical or non numerical. Numerical affinity represents the sum of the frequencies of queries which access simultaneously P_i and P_j . Non numerical affinity underlines the implication relation between predicates P_i and P_j

Let $P = \{P_i, i = 0 \text{ to } 8\}$ be the set of predicates used by Q.

 $\begin{array}{l} \textbf{P_0: ObNature} = ``vehicule''\\ \textbf{P_1: DC Similar(0.3) ((12; 10; 13), (14; 15; 16), (20; 20; 20))\\ \textbf{P_2: ObColor} = ``rted''\\ \textbf{P_3: DC Similar(0.2) ((12; 10; 13), (14; 15; 16), (20; 20; 20))\\ \textbf{P_4: ObNature} = ``truck''\\ \textbf{P_5: DC Similar(0.1) ((9; 8; 7), (7; 8; 7), (10; 11; 10))\\ \textbf{P_6: DC Similar(3) ((12; 10; 13), (14; 15; 16), (20; 20; 20))\\ \textbf{P_7: DC Similar(1) ((12; 10; 13), (14; 15; 16), (20; 20; 20))\\ \textbf{P_8: DC Similar(1) ((9; 8; 7), (7; 8; 7), (10; 11; 10))\\ \end{array}$

P contains traditional predicates (P_0 , P_2) as well as multimedia predicates (P_1 , P_3 , P_4 , P_5 , P_6 , P_7 , P_8). Note P_1 , P_3 and P_5 are range query predicates (the number between brackets being a real value – similarity range ε), while P_6 , P_7 and P_8 are KNN predicates (the number between brackets being an integer value – number of objects k to be returned by the predicate). Also note that *DC* represents a composite *Dominant Color* feature vector stating the three consecutive dominant colors in an image, in RGB color space. For example, *DC*₁ of predicate p_1 underlines dominant colors C(12; 10; 13), C'(14; 15; 16) and C''(20; 20; 20).

By reading the updated PAM, one can clearly point out the multimedia implication rules defined in the paper:

- Predicate P_3 ($\varepsilon_3 = 0.2$, $V_3 = ((12; 10; 13), (14; 15; 16), (20; 20; 20)))$ implies P_1 ($\varepsilon_1 = 0.3$, $V_1 = ((12; 10; 13), (14; 15; 16), (20; 20; 20)))$ having:
 - $V_1 = V_3$ and $\varepsilon_1 > \varepsilon_3$
- Predicate $P_5 (\epsilon_5 = 0.1 \text{ "max"}, V_5 = ((9; 8; 7), (7; 8; 7), (10; 11; 10)))$ implies $P_1 (\epsilon_1 = 0.3, V_1 = ((12; 10; 13), (14; 15; 16), (20; 20; 20)))$ having:
 - $\varepsilon_1 > \varepsilon_5$, $dist(V_1, V_5) \le \varepsilon_1 \varepsilon_5$
- No implication can be identified between predicates P₃ and P₅ having:
- $dist(V_3, V_5) > \varepsilon_3 \varepsilon_5$ (similarity circle intersection/exclusion)
- Predicate $P_7 (k_7 = 1, V_7 = ((12; 10; 13), (14; 15; 16), (20; 20; 20)))$ implies predicate $P_6 (k_6 = 3, V_6 = ((12; 10; 13), (14; 15; 16), (20; 20; 20)))$ having:
 - $V_6 = V_7$ and $k_6 > k_7$
- No implication can be identified between P₆ (orP₇) and P₈, having:
 V₈ ≠ V₆ (correspondingly V₇)

	PU	M and	I PAM											
F	ile Q	uery N	lodels Ir	troduction to Fragmentation										
User queries and corresponding access frequencies mmaized query string		Prec Q(0) Q(1) Q(2) Q(3) Q(4) Q(5)	icate Us P(0 : 1 : 1 : 0 : 1 : 1 : 0	iage Ma) P(1) 1 0 0 0 0 0	trix P(2) P(3) 0 0 1 1 1 0 0 0 1 0 1 0	P(4) P(5) 0 0 1 1 0 0 0 0 1 0	P(6) P(7) 0 0 0 0 1 0 1 0 0 1 0 0	P(8) F 0 0 0 0 1	req 5 10 5 5 5 5 5 5		PUM			
	•	2 3 4 5	10 5 5 5	Select sr1.* From tbl_StillR Select sr1.* From tbl_StillR Select sr1.* From tbl_StillR Select sr1.* From tbl_StillR			Γ	Pı	edicat	te Usa	ige N	/latrix	:	
	1		1	•	Prov									MAP
Applied Predicates F0(1): ODKature = "vehicla" P(1): DDKature = "vehicla" P(2): ODColor = "ed" P(3): DDCSmiller(2): 12:10:13:14:15:16:20:20:20 P(4): ODKature = "true" P(5): DDCSmiller(2): 12:10:13:14:15:16:20:20:20 P(4): ODKature = "true" P(5): DDCSmill(1): 39:77:87:10:11:10 P(6): DDCSmill(1): 32:77:87:10:11:10 P(6): DDCSmill(1): 32:77:87:10:11:10 P(6): DDCSmill(1): 32:77:87:10:11:10 P(6): DDCSmill(1): 32:77:87:10:11:10				P(0) P(1) P(2) P(3) P(4) P(5) P(6) P(7) P(8)	P(0) 20 5 10 5 0 5 5 0	P(ttrix 1) P(2 1) 1 1) 2 1) 2 1) 2 1) 1 1) 1 1) 0 1) 5 1) 5 11 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1) P(3) 0 5 5 5 5 0 0 0 0 0 0	P(4) 0 15 0 15 10 0 5	P(5) 0 <= 10 0 10 10 0 0 0	P(6) 5 0 0 0 0 5 => 0	P(7) 5 0 5 0 0 5 0 0 5 0	P(8) 0 0 5 0 5 0 0 5 0 0 5	
Predicates invoked in user queries				Updated Predicate Affinity Matrix										

Fig. 6. Updated Predicate Affinity Matrix.

Disregarding our multimedia implication rules would yield, in the present example, a PAM with only numerical affinities.

The PUM and uPAM make up the inputs to the NHP primary horizontal partitioning algorithm [11, 2], not being implemented yet in our prototype.

5.2 Timing Analysis

We have shown that the complexity of our physical similarity implication simplifies to $O(n_v \times n_p^2)$. We verified the formula experimentally, the timing results being presented in *Fig. 7*.



Fig. 7. Timing results

The experiment was carried out on a Pentium 4 PC (2.8 Ghz CPU, 798 Mhz bus, 512 MB RAM). One can see that the time to compute similarity implications grows in a polynomial (quadratic) fashion with the number of predicates involved. Our experiments also show that feature vector dimension affects time complexity, owing to predicate distance computations (weighted Euclidian distance).

6 Conclusion and Future Work

In this paper, we proposed an approach for the Primary Horizontal Fragmentation of multimedia databases, by extending existing fragmentation methods. Following the definition of a multimedia type, we identified the need to classify multimedia objects corresponding to the same type, in order to achieve consistent horizontal fragmentation criteria. The "Type Fragmentation" phase could be then followed by the PHF of each generated type. The original idea of emerging new multimedia operators allowed the adaptation of existing fragmentation procedures to partition multimedia data. We concentrated our efforts on the primary horizontal fragmentation of unstructured multimedia data, emphasizing the impact of multimedia predicate implications in optimizing multimedia fragments.

Future directions include the introduction of semantic-based multimedia predicates. Our future goals also incorporate generating a multimedia conceptual schema, including the derived horizontal fragmentation process, and optimizing, if possible, the used fragmentation methods (semantic implication is yet to be developed). Likewise, multimedia vertical fragmentation and XML fragmentation will be talked in upcoming studies.

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